

# IR-CGR: Time-Aware Inter-Region Contact Graph Routing for Scalable Space DTNs

---

Shota Suzuki (Keio University, Japan)

Juan A. Fraire (Inria, INSA Lyon, France; CONICET, Argentina)

Keisuke Uehara (Keio University, Japan)

@ STINT Workshop, IEEE WiSEE 2026

# Outline

## 1. Background and existing methods

CGR has a scalability problem; LRK loses timing information.

## 2. Proposal: IR-CGR

A time-aware inter-region routing method.

## 3. Evaluation Setup

Two benchmarks for robustness and scalability.

## 4. Orbital Benchmark

Delivery robustness in a realistic lunar scenario.

## 5. Scaling Benchmark

Computational scaling with network size.

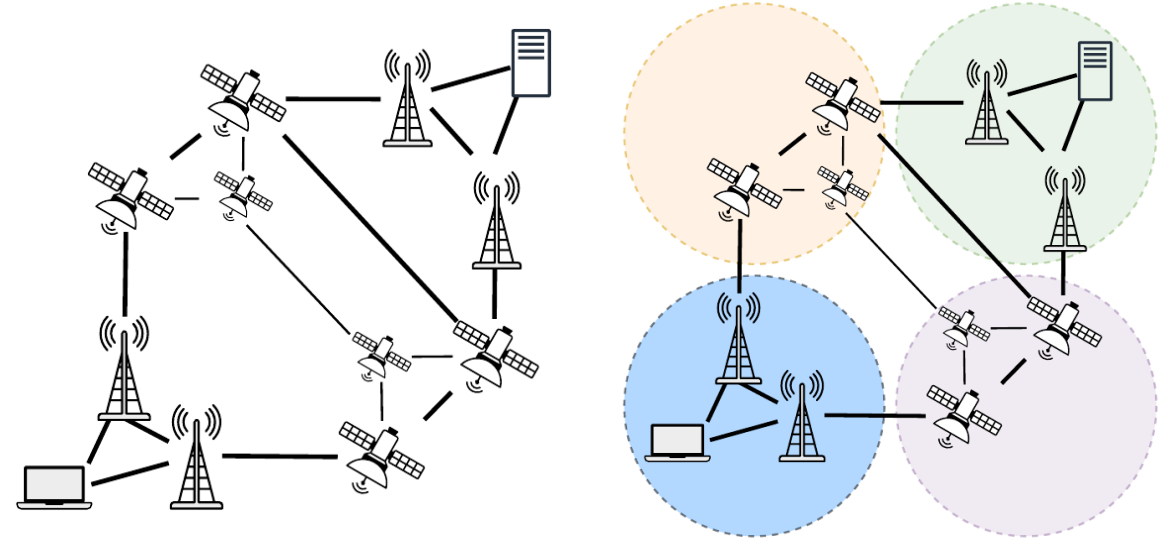
## 6. Summary

Benefits and trade-offs of IR-CGR.

# Background: the CGR scalability problem

## CHALLENGE

- DTN protocols are being standardized and adopted for future space networking.
- CGR exploits predicted future contacts to compute routes in scheduled DTNs.
- Its worst-case route-search complexity is  $O(|C|^2)$  ( $\times |C|$  = number of contacts)



Space Network without regions:  
every node searches the full plan

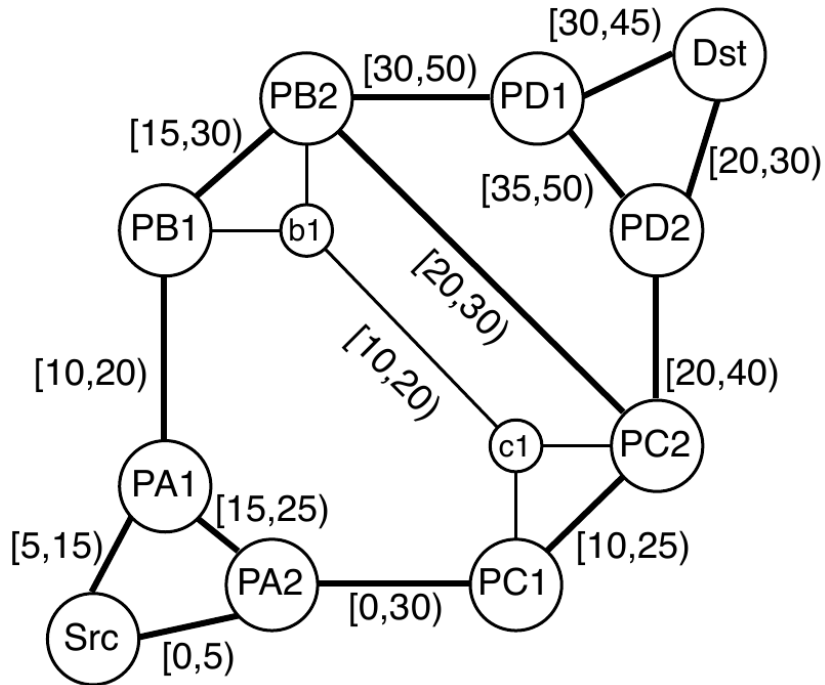
Space Network with regions:  
each node searches its region

**Contact plans grow with network size, while onboard resources do not.**

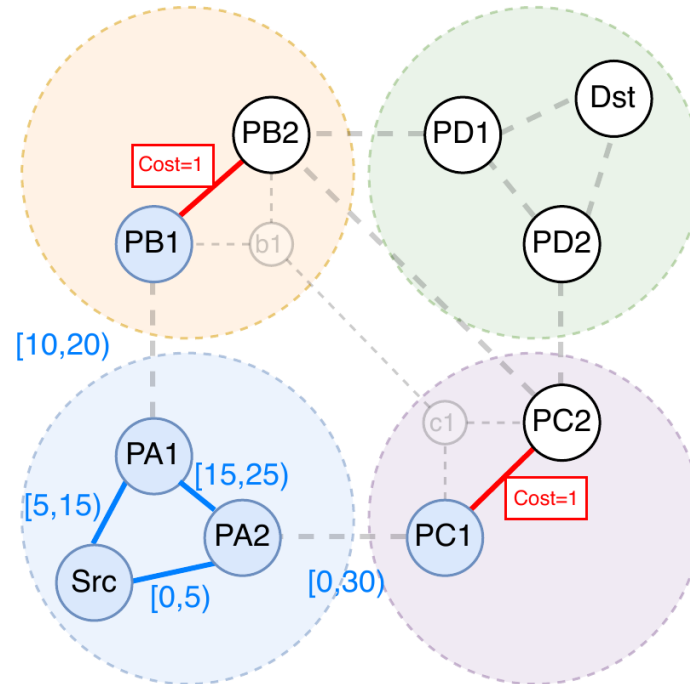
**➡ Each node's knowledge of the network must be bounded.**

# Existing method: Local Regionalized Knowledge (LRK) [2]

**Global Full Knowledge (GFK)**  
(= Conventional CGR without Region)



**Local Regionalized Knowledge (LRK)**



## Intra-region

- Regional CGR on the regional contact plan
- Contact windows preserved  
→ **time-aware**

## Inter-region

- Static passageway graph with fixed crossing costs
- Contact windows not represented → **time-static**

**LRK keeps each node's routing knowledge small.  
But can it work well when timing information is lost at region crossings?**

# Outline

## 1. Background and existing methods

CGR has a scalability problem; LRK loses timing information.

## 2. Proposal: IR-CGR

A time-aware inter-region routing method.

## 3. Evaluation Setup

Two benchmarks for robustness and scalability.

## 4. Orbital Benchmark

Delivery robustness in a realistic lunar scenario.

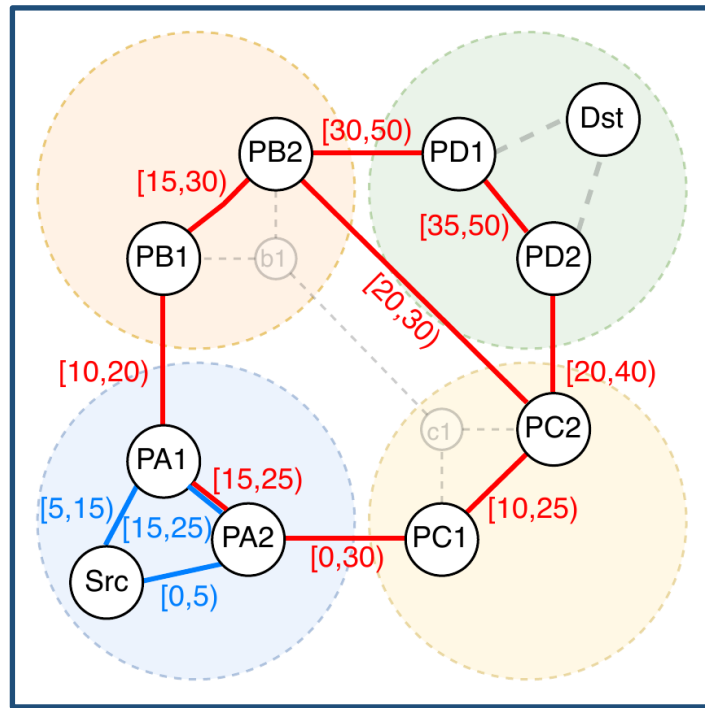
## 5. Scaling Benchmark

Computational scaling with network size.

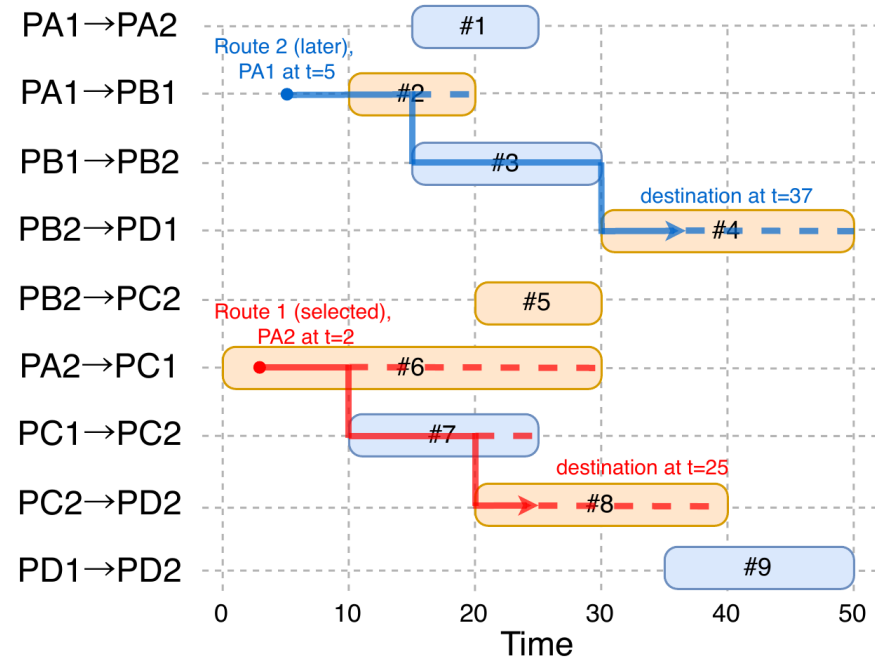
## 6. Summary

Benefits and trade-offs of IR-CGR.

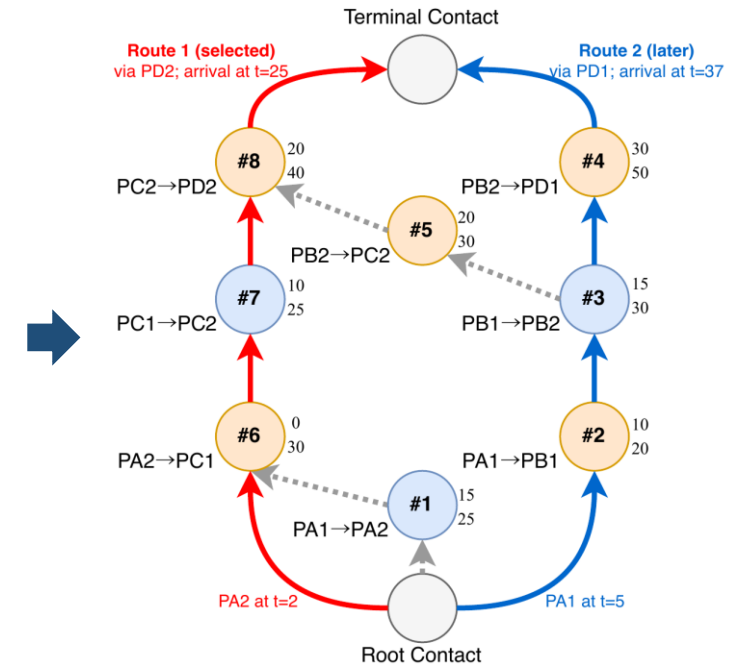
# Proposal: Inter-Region Contact Graph Routing (IR-CGR)



Divide network into regions



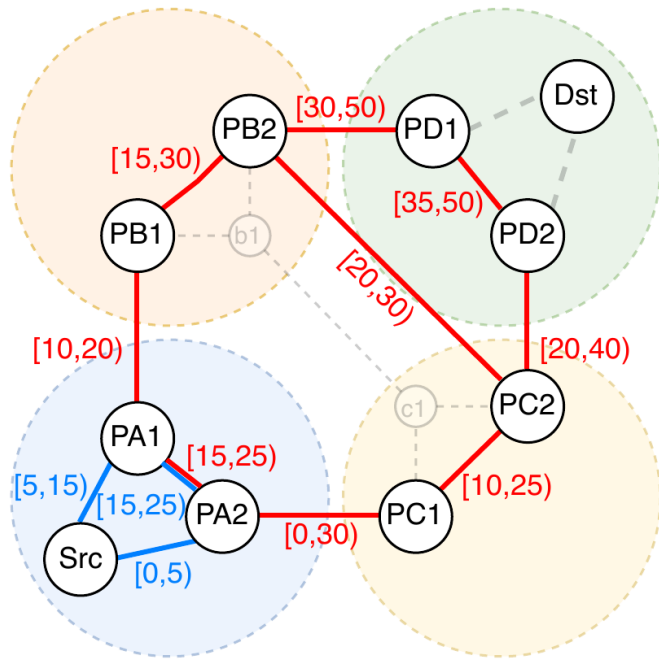
IR-Contact Plan (IRCP)



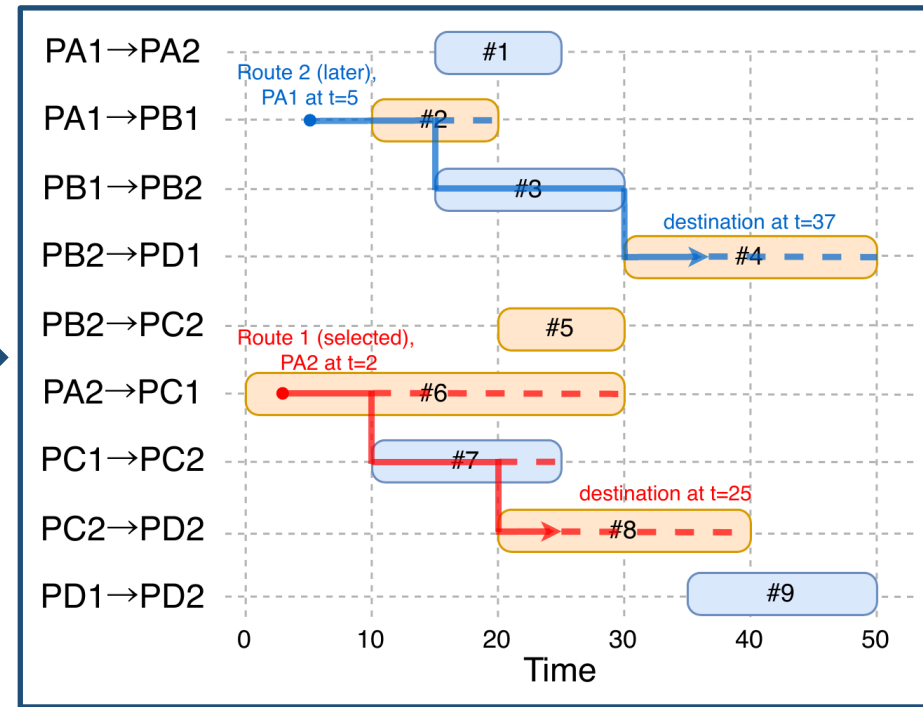
IR-Dijkstra

- Same regionalization as LRK
- But inter-region crossings keep their contact windows

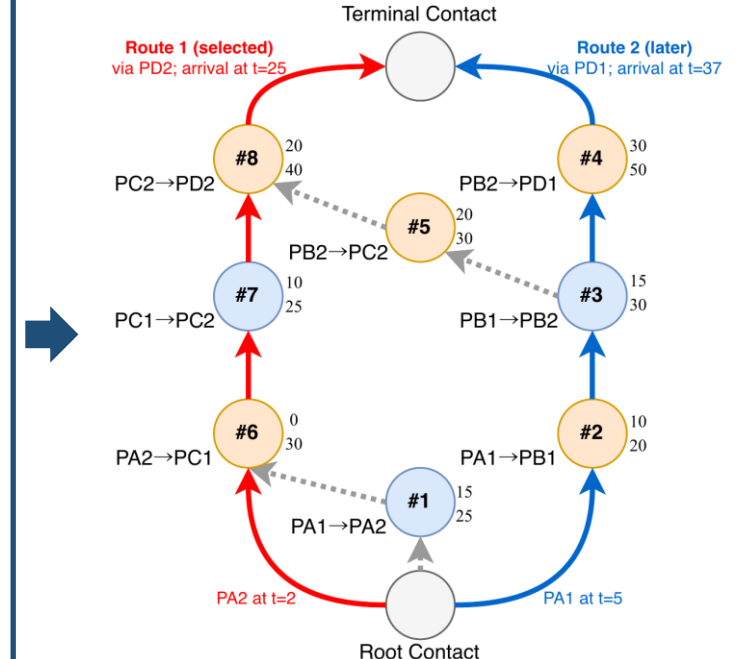
# Proposal: Inter-Region Contact Graph Routing (IR-CGR)



Divide network into regions



IR-Contact Plan (IRCP)

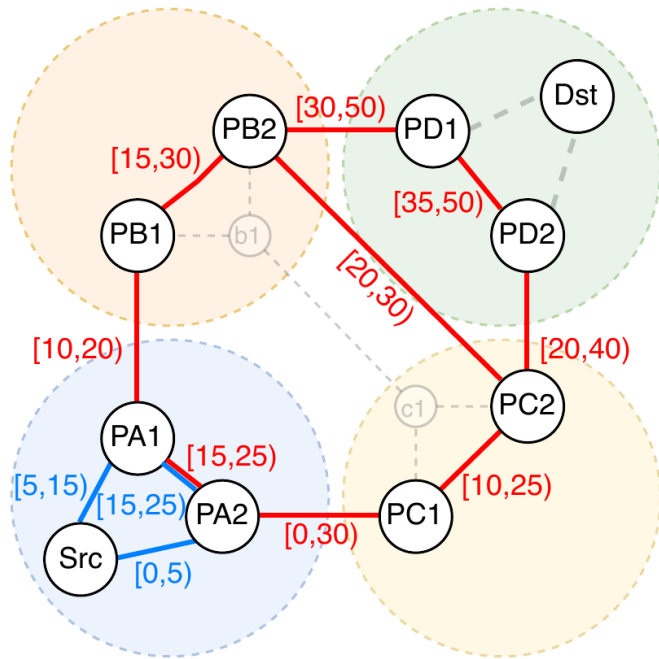


IR-Dijkstra

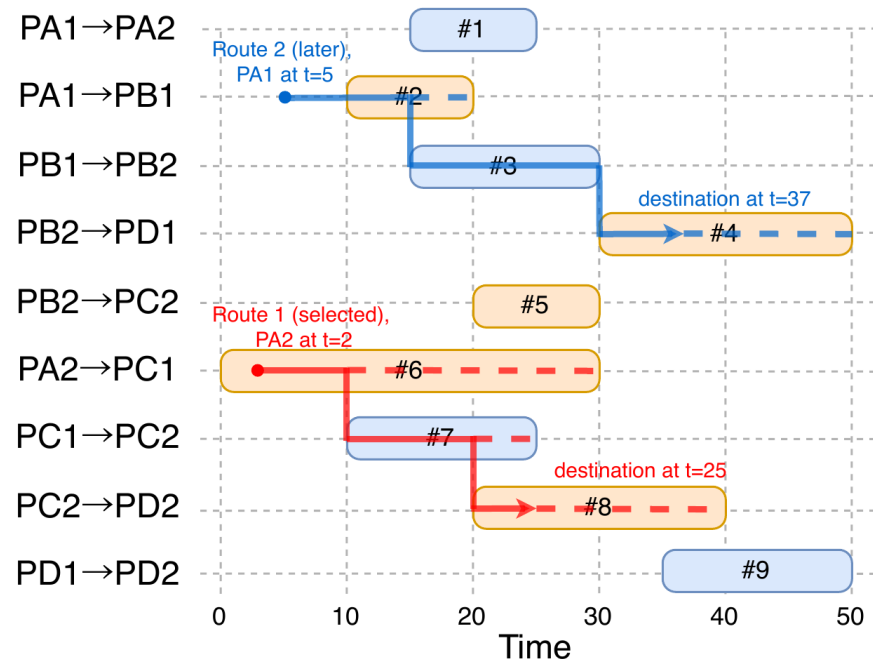
- **INTER IR-Contact** : actual cross-region contacts
- **INTRA IR-Contact** : sampled passageway-to-passageway transit
- IRCP is built on the ground for each planning horizon

✂Planning horizon = the time span covered by the contact plan

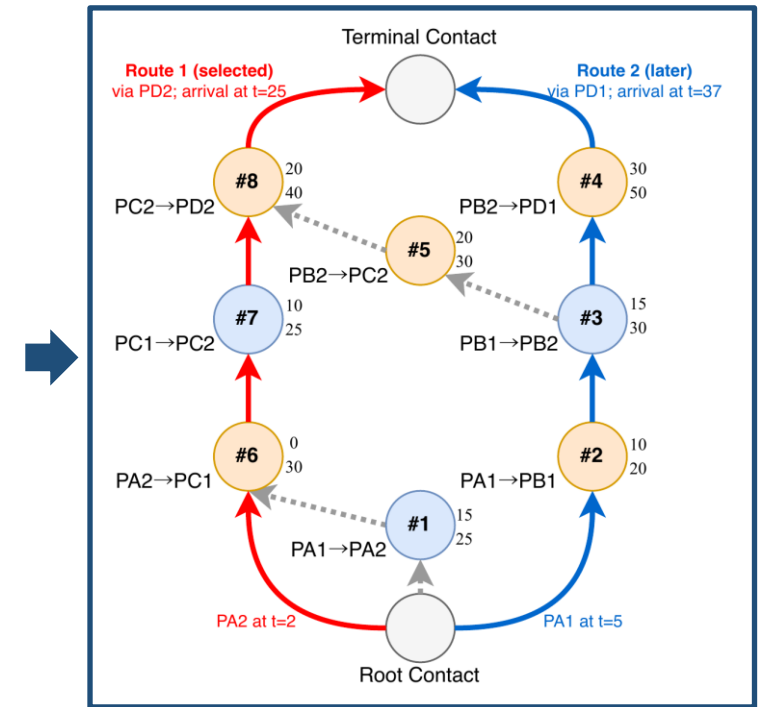
# Proposal: Inter-Region Contact Graph Routing (IR-CGR)



Divide network into regions



IR-Contact Plan (IRCP)



IR-Dijkstra

- IR-Dijkstra searches the IRCP for the earliest-arrival passageway sequence
- Regional CGR forwards the bundle to the next passageway

**Same time-aware routing idea as CGR, applied one level higher.**

# Outline

## 1. Background and existing methods

CGR has a scalability problem; LRK loses timing information.

## 2. Proposal: IR-CGR

A time-aware inter-region routing method.

## 3. Evaluation Setup

Two benchmarks for robustness and scalability.

## 4. Orbital Benchmark

Delivery robustness in a realistic lunar scenario.

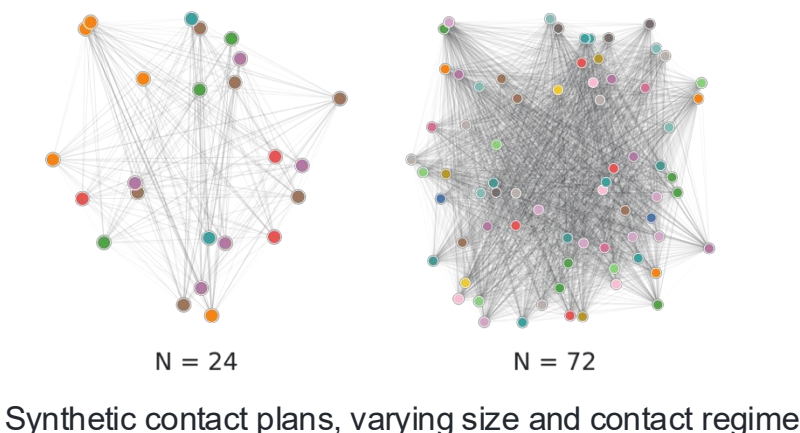
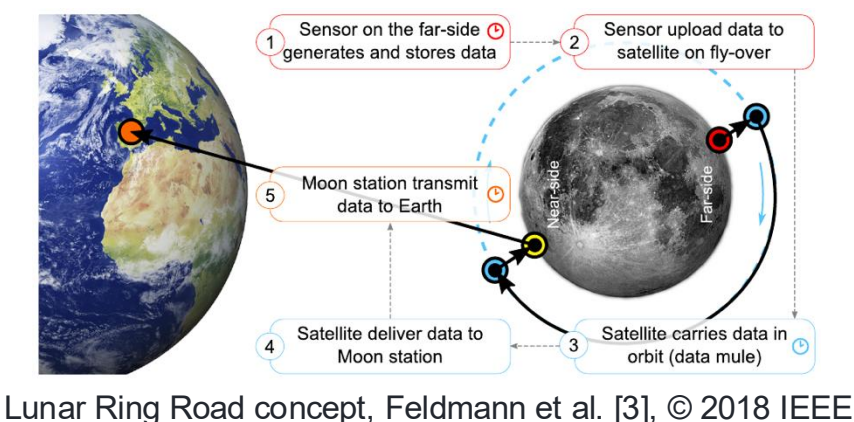
## 5. Scaling Benchmark

Computational scaling with network size.

## 6. Summary

Benefits and trade-offs of IR-CGR.

# Evaluation setup



## Orbital Benchmark (Lunar Ring Road [3])

## Scaling Benchmark (DDRCM)

|                    |                                                                    |                                                                |
|--------------------|--------------------------------------------------------------------|----------------------------------------------------------------|
| Network            | 23 nodes, orbital contact plan                                     | 24 to 72 nodes, synthetic plan                                 |
| Links              | with/without inter-satellite links (ISLs) between lunar satellites | three distance-dependent contact regimes ( $\beta = 1, 2, 4$ ) |
| Sweep              | 4 to 12 regions                                                    | 4, 8 or 12 nodes per region                                    |
| Seeds              | 100                                                                | 20                                                             |
| Compared / metrics | GFK, LRK, IR-CGR; delivery ratio, delay, cost                      |                                                                |

**Orbital Benchmark: Evaluates delivery robustness in a realistic orbital scenario**  
**Scaling Benchmark: Evaluates computational scaling with network size**

# Outline

## 1. Background and existing methods

CGR has a scalability problem; LRK loses timing information.

## 2. Proposal: IR-CGR

A time-aware inter-region routing method.

## 3. Evaluation Setup

Two benchmarks for robustness and scalability.

## 4. Orbital Benchmark

Delivery robustness in a realistic lunar scenario.

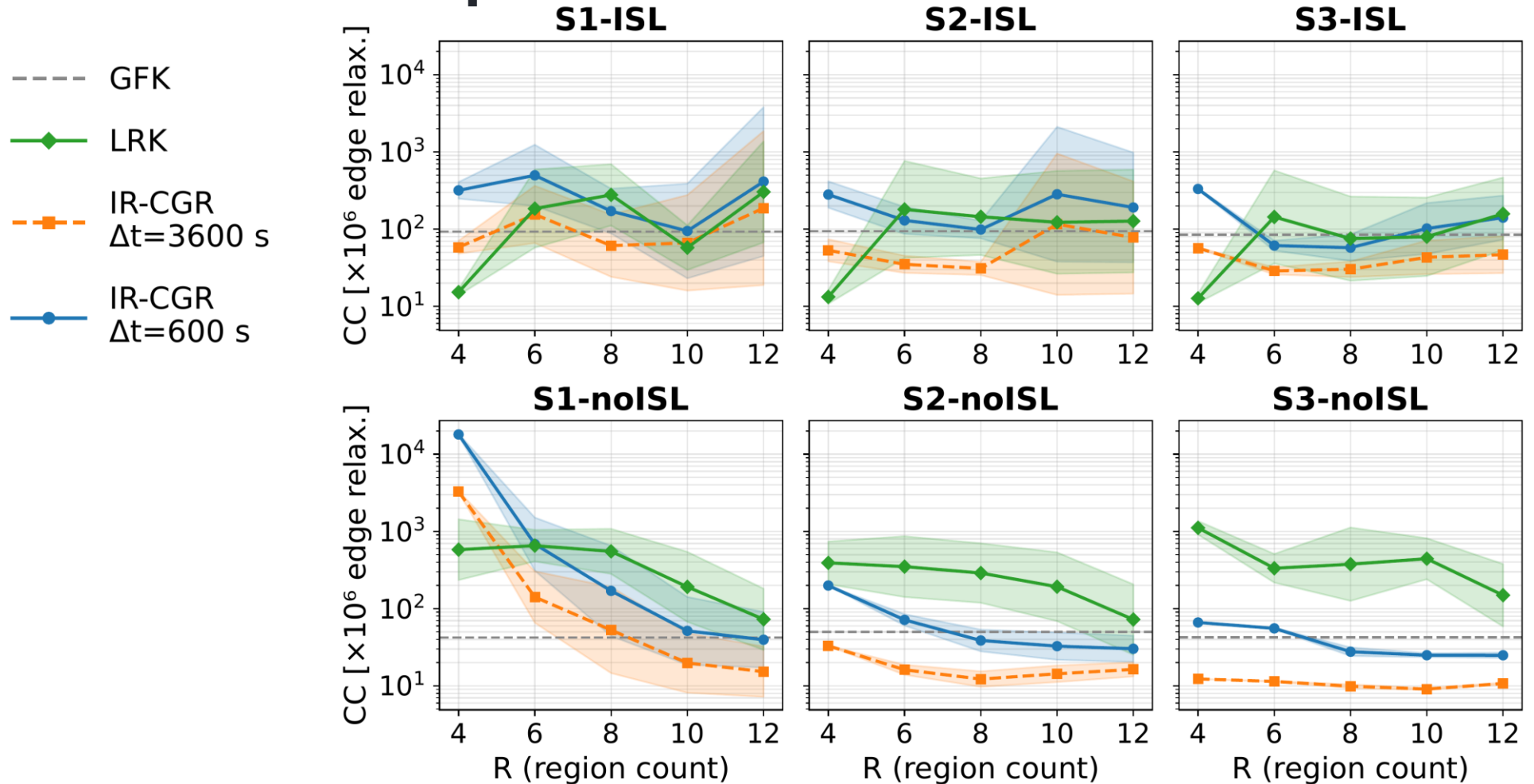
## 5. Scaling Benchmark

Computational scaling with network size.

## 6. Summary

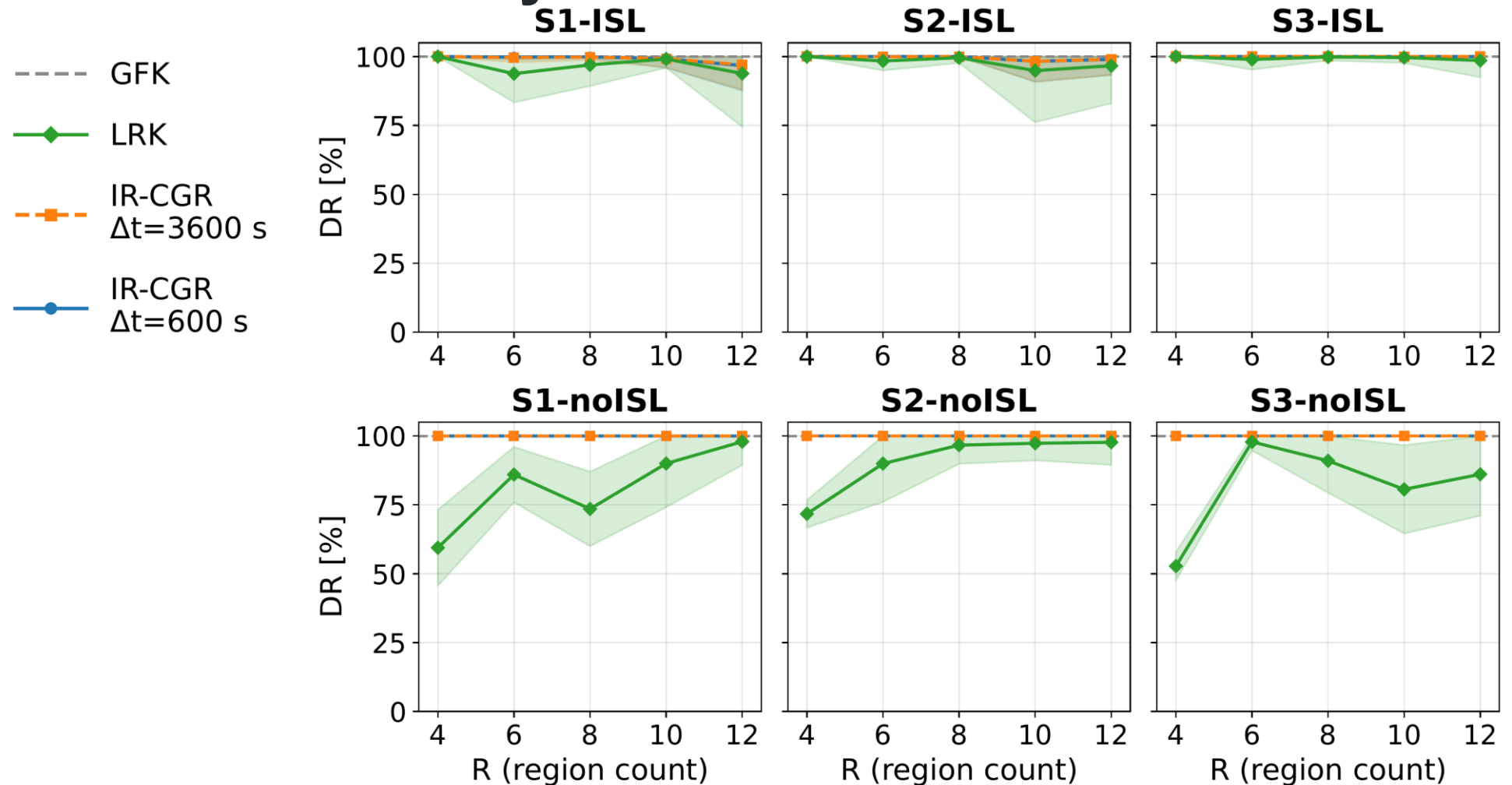
Benefits and trade-offs of IR-CGR.

# Orbital Benchmark: Computational Cost



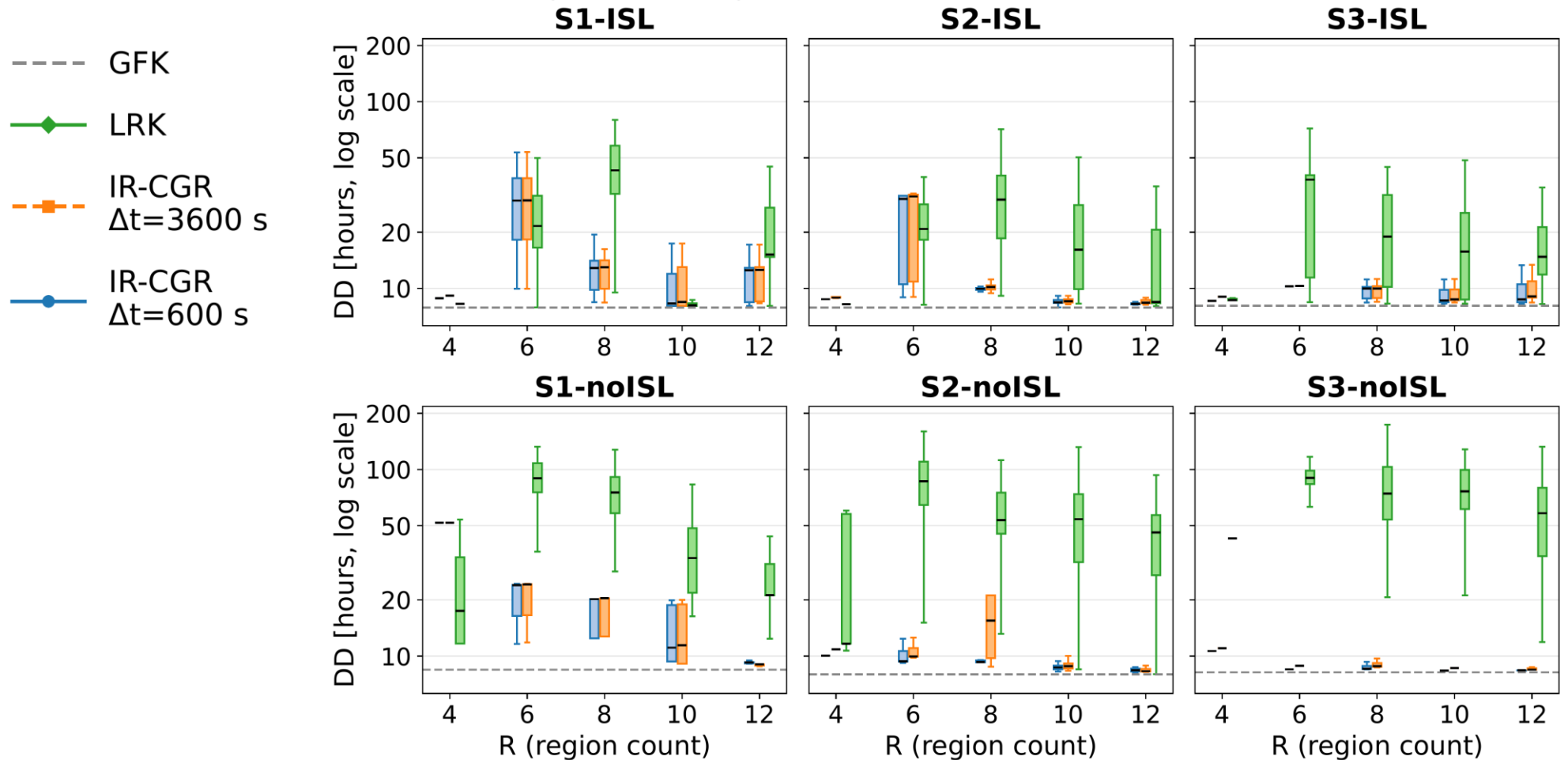
- With coarse partitioning, both IR-CGR and LRK can cost more than GFK.
  - With suitable partitioning, IR-CGR can cost less than both GFK and LRK.
- The operator's partitioning choice therefore determines the routing-cost trade-off.

# Orbital Benchmark: Delivery Ratio



- LRK loses delivery when inter-region connectivity becomes sparse
- IR-CGR ( $\Delta t=3600$  s) maintains full delivery across the evaluated no-ISL cases.

# Orbital Benchmark: Delivery Delay



- With suitable partitioning, IR-CGR stays close to GFK; coarse partitions can increase its delay.
- LRK shows much larger and more variable delays, especially without ISLs.

# Outline

## 1. Background and existing methods

CGR has a scalability problem; LRK loses timing information.

## 2. Proposal: IR-CGR

A time-aware inter-region routing method.

## 3. Evaluation Setup

Two benchmarks for robustness and scalability.

## 4. Orbital Benchmark

Delivery robustness in a realistic lunar scenario.

## 5. Scaling Benchmark

Computational scaling with network size.

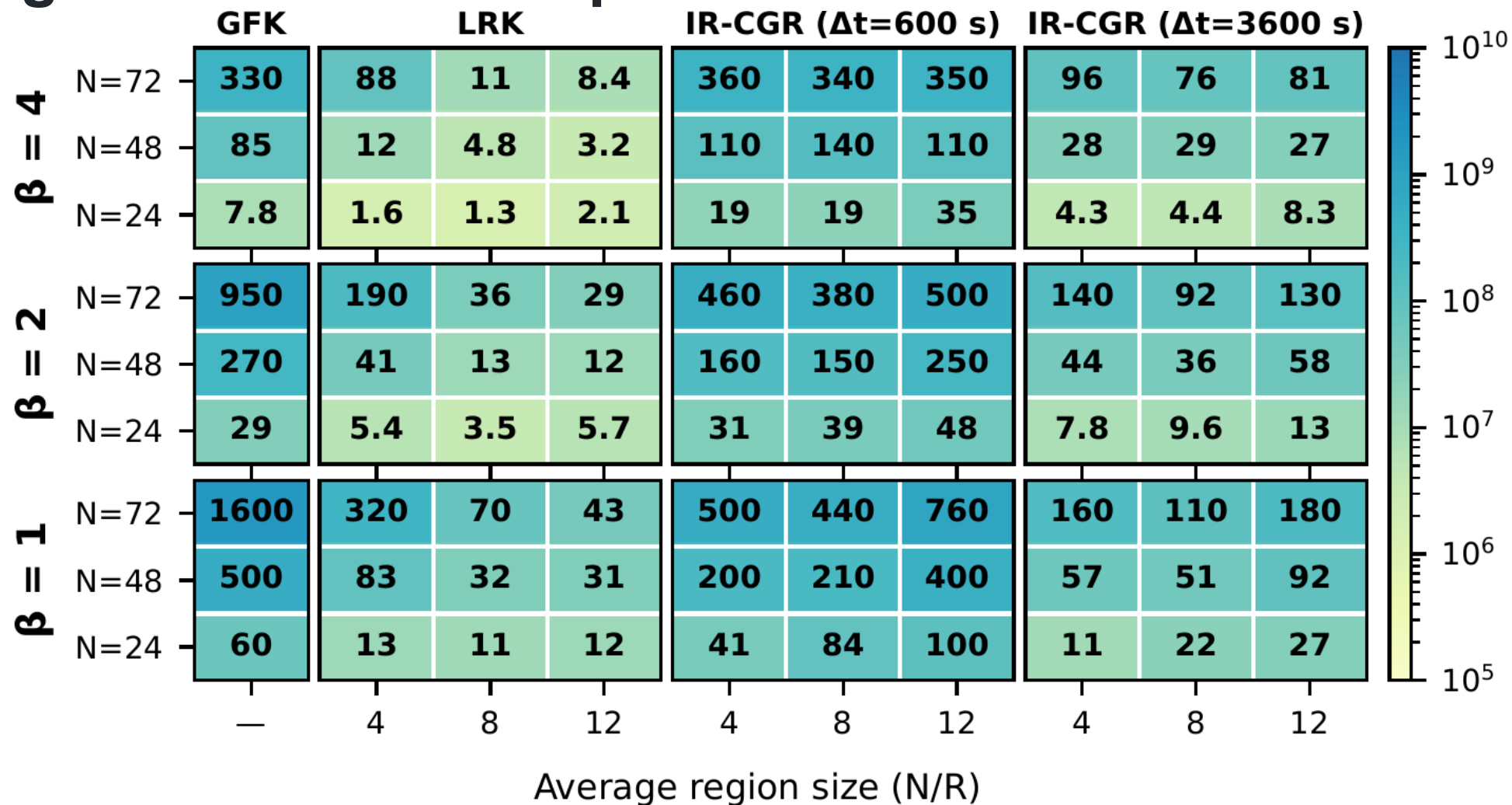
## 6. Summary

Benefits and trade-offs of IR-CGR.

# Scaling Benchmark: Computational Cost

Every design delivers every bundle on this grid.

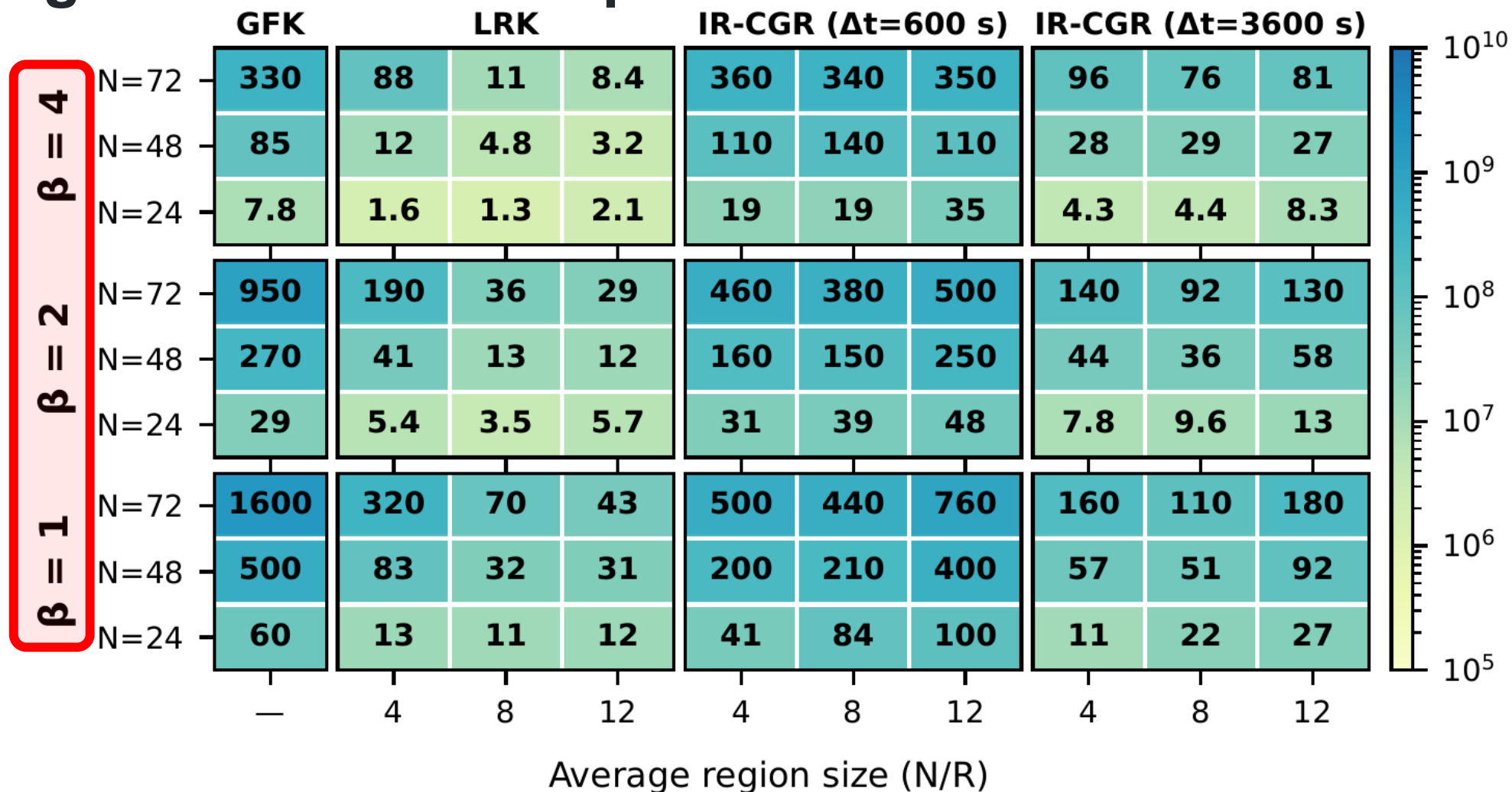
Cell = 20-seed mean; CC = edge relaxations



# Scaling Benchmark: Computational Cost

Every design delivers every bundle on this grid.

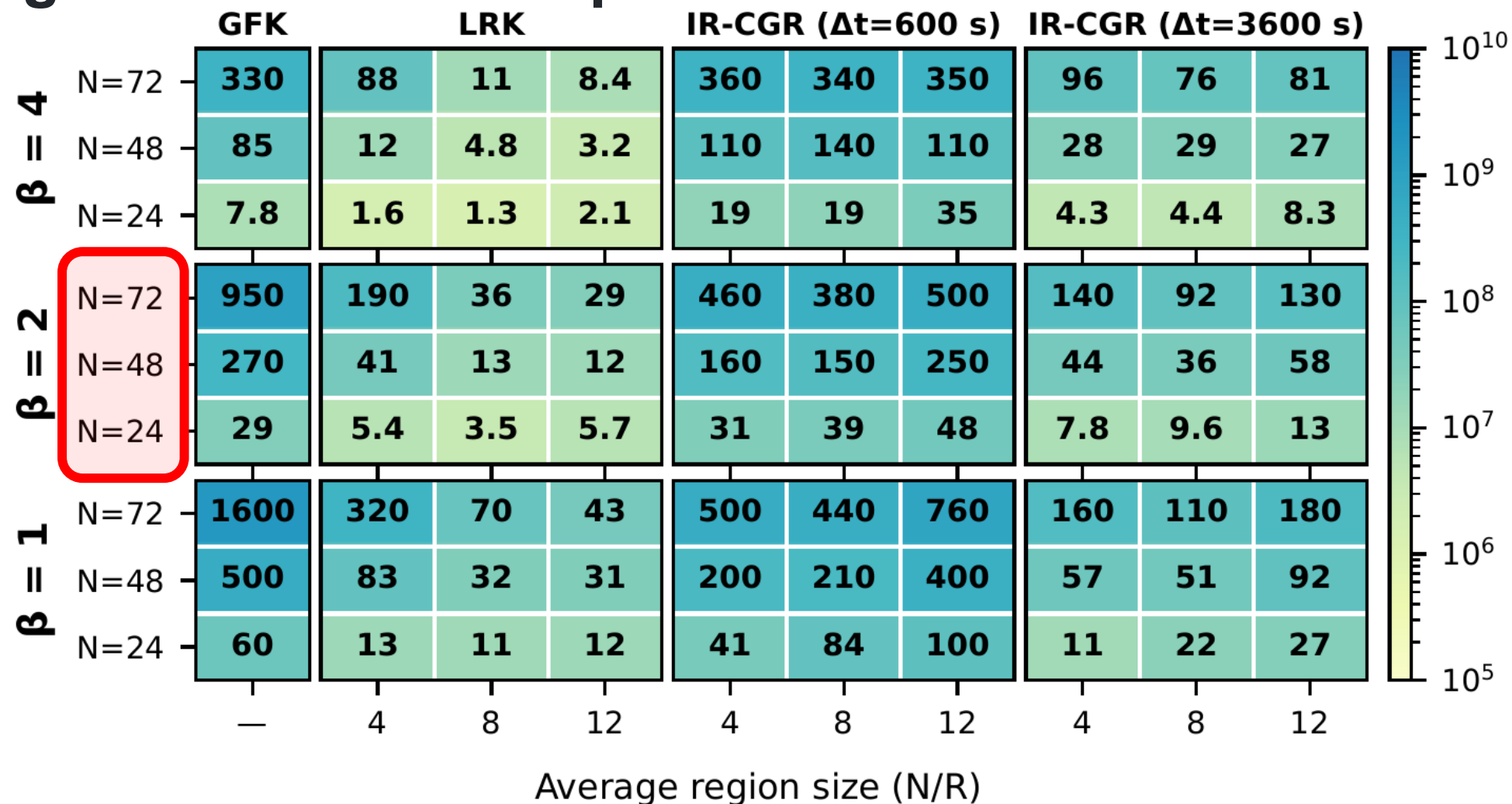
Cell = 20-seed mean; CC = edge relaxations



# Scaling Benchmark: Computational Cost

Every design delivers every bundle on this grid.

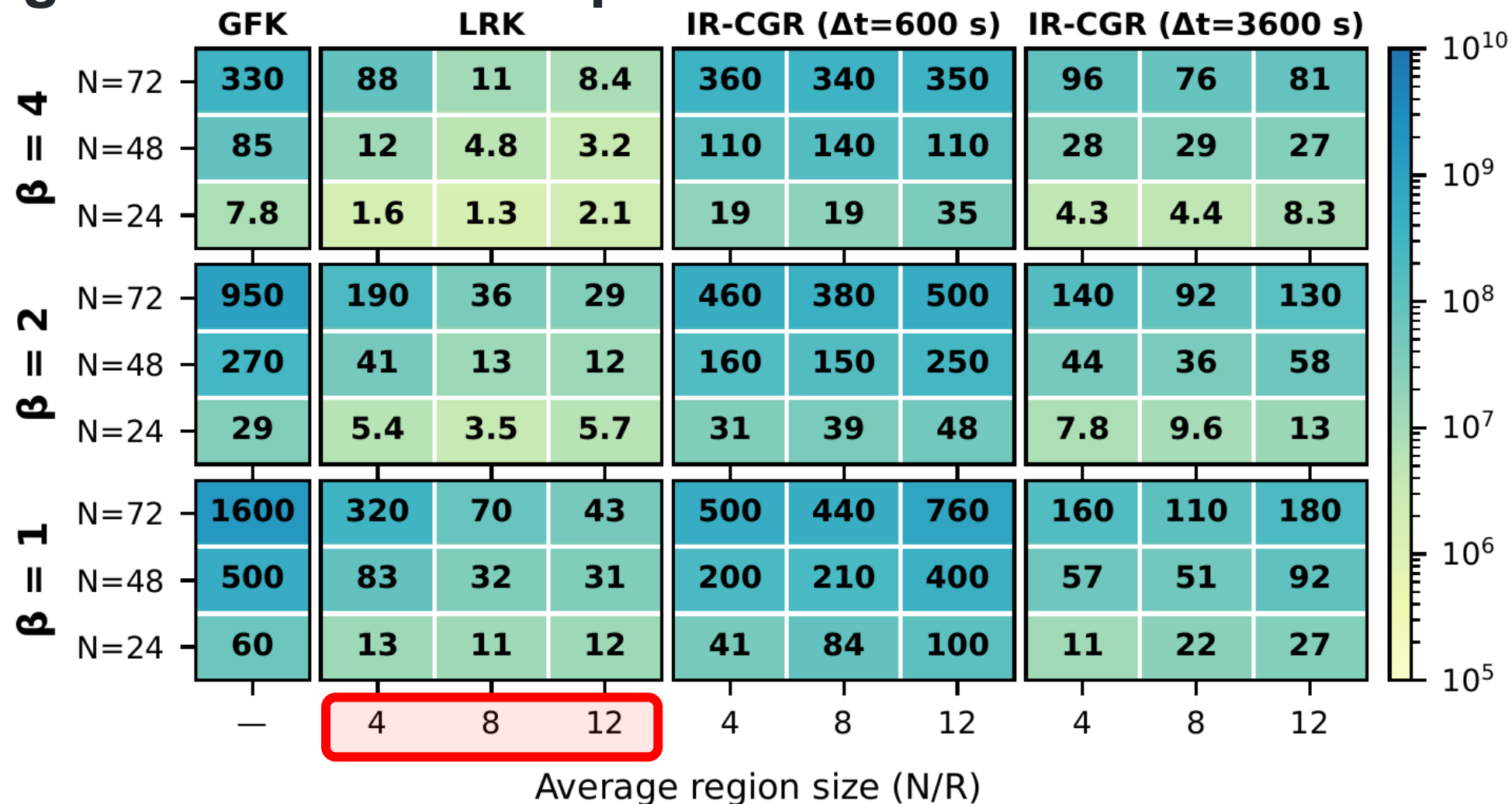
Cell = 20-seed mean; CC = edge relaxations



# Scaling Benchmark: Computational Cost

Every design delivers every bundle on this grid.

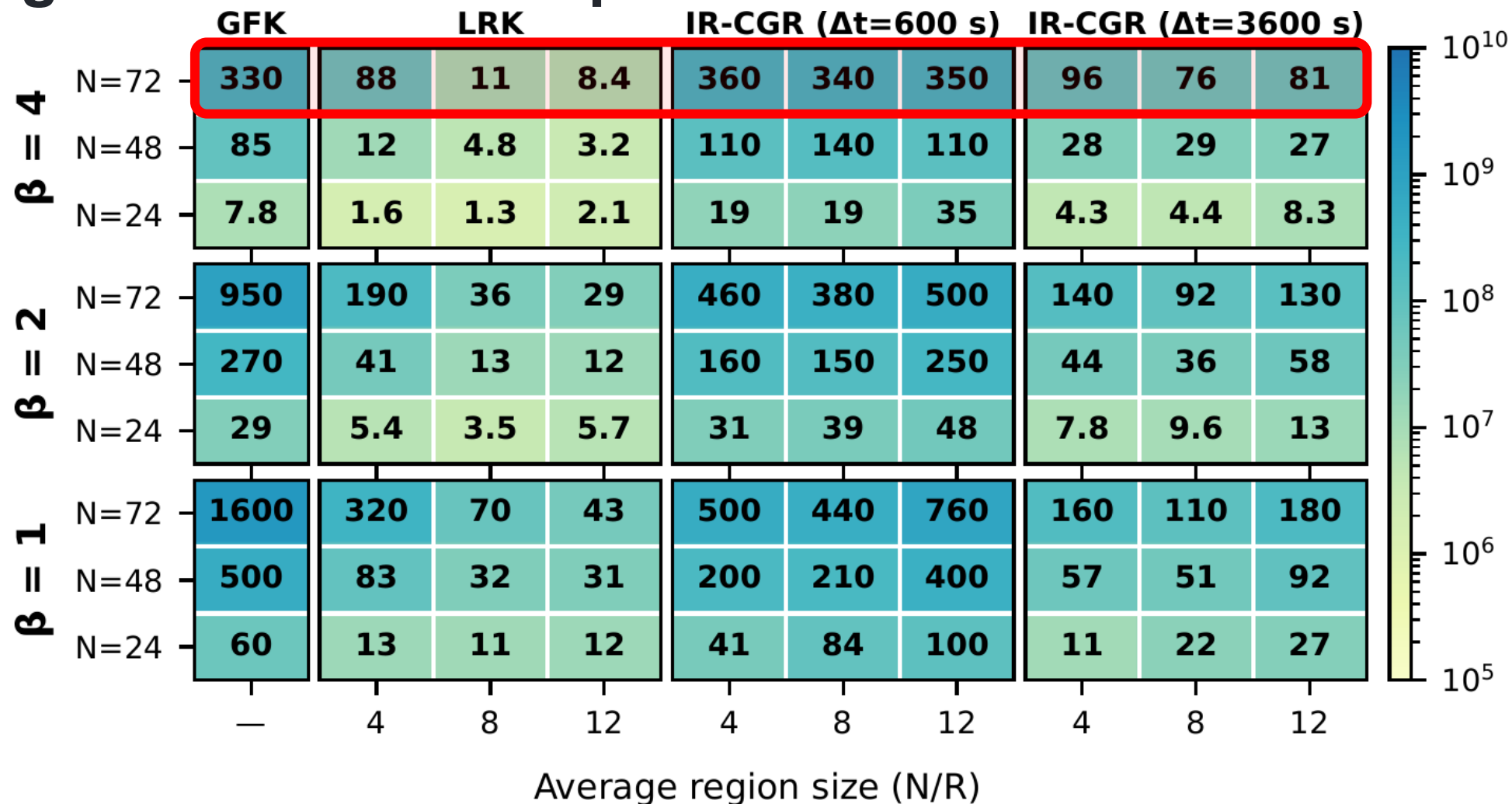
Cell = 20-seed mean; CC = edge relaxations



# Scaling Benchmark: Computational Cost

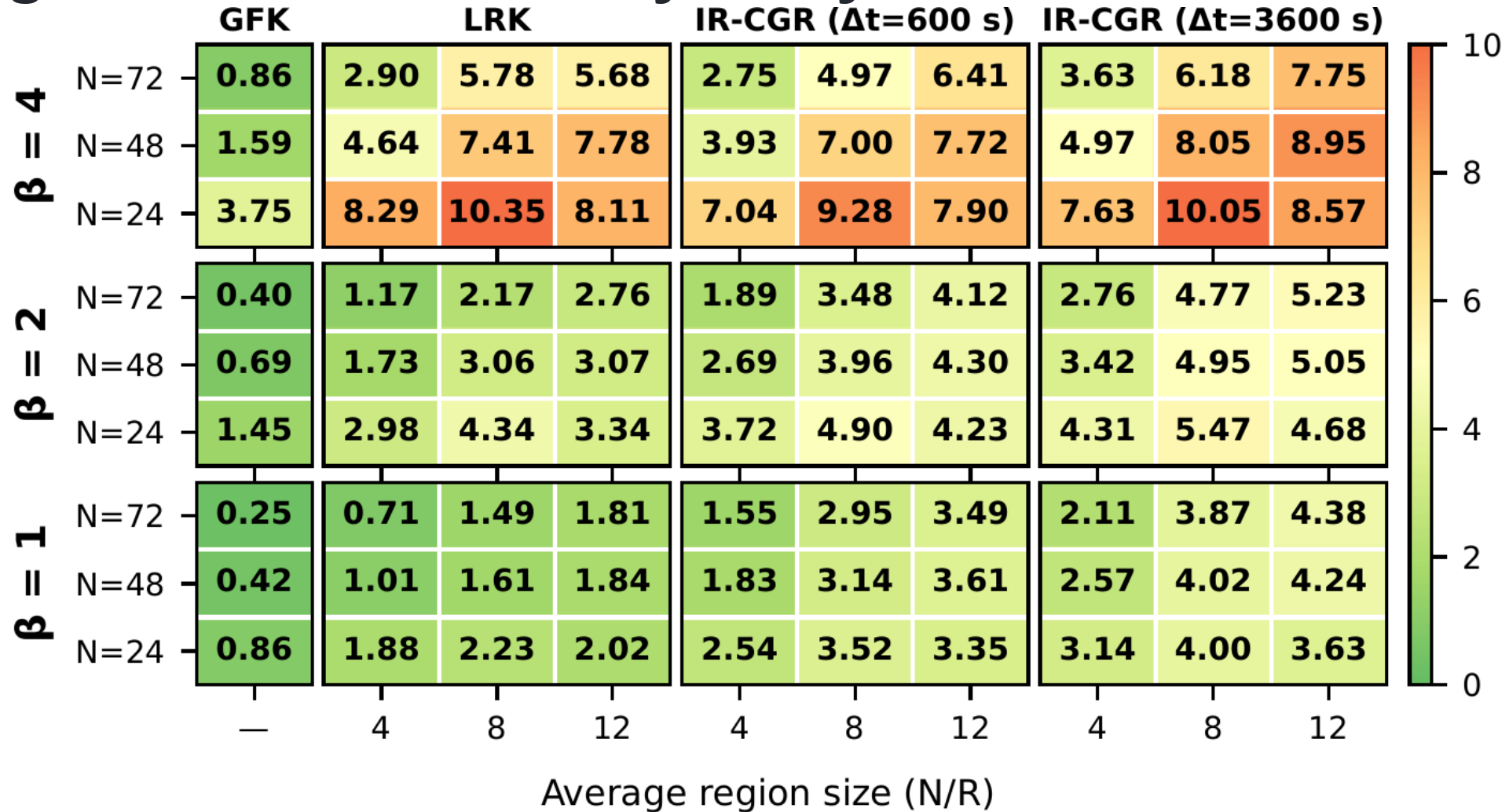
Every design delivers every bundle on this grid.

Cell = 20-seed mean; CC = edge relaxations



IR-CGR is usually more expensive than LRK,  
but can remain much cheaper than full-plan GFK as the network grows.

# Scaling Benchmark: Delivery Delay



Both regionalized methods show substantially higher delay than GFK.  
In this benchmark, IR-CGR often does not improve delay over LRK.

# Outline

## 1. Background and existing methods

CGR has a scalability problem; LRK loses timing information.

## 2. Proposal: IR-CGR

A time-aware inter-region routing method.

## 3. Evaluation Setup

Two benchmarks for robustness and scalability.

## 4. Orbital Benchmark

Delivery robustness in a realistic lunar scenario.

## 5. Scaling Benchmark

Computational scaling with network size.

## 6. Summary

Benefits and trade-offs of IR-CGR.

# Summary

## Problem

LRK reduces routing knowledge, but its static inter-region routing loses contact timing information.

## Proposal

IR-CGR keeps contact windows between regions and applies CGR-style earliest-arrival routing on a compact IRCP.

## Result

Overall, IR-CGR is most valuable under sparse, time-varying inter-region connectivity, while LRK remains attractive when inter-region connectivity is sufficiently rich and stable.

# References

- [1] J. A. Fraire, O. De Jonckère and S. C. Burleigh, “Routing in the Space Internet: A Contact Graph Routing Tutorial,” *Journal of Network and Computer Applications*, vol. 174, p. 102884, 2021.
- [2] P. G. Madoery, J. A. Fraire, F. D. Raverta, J. M. Finochietto and S. C. Burleigh, “Managing Routing Scalability in Space DTNs,” in *Proc. 6th IEEE Int. Conf. Wireless for Space and Extreme Environments (WiSEE)*, Huntsville, AL, USA, 2018, pp. 177–182.
- [3] M. Feldmann, J. A. Fraire and F. Walter, “Tracking Lunar Ring Road Communication,” in *Proc. IEEE Int. Conf. Communications (ICC)*, Kansas City, MO, USA, 2018, pp. 1–7